



LOGIC AND COMPUTER SCIENCE 31 P

EDITED BY P. ODIFREDDI

Two Extensions of Curry's Type Inference System by F. Cardone and M. Coppo

MARIANGIOLA

photo by J.W. Klop 1978

NY

ETAPS'02 ???

 \cap

 \mathcal{M}

JARIANGIOLA

photo by J.W. Klop 1978

ETAPS'02 ???

Mini



 \bigcirc

0

 \bigcirc

NY





Session Types in a Nutshell

SESSION = STRUCTURED SEQUENCE OF COMMUNICATION

send(int).send(int).receive(bool)

"...Session Types *structure* a *series of interactions* in a simple and concise syntax and ensure *type safe communication*."

A Protocol





- Protocol: Buyer-Seller
- Description: Alice buying a book

send(string).receive(int).⊕{ok: send(string).receive(date), quit:end}
receive(string).send(int).&{ok: receive(string).send(date), quit: end}



Are we compatible?

send(int).send(int).receive(bool)



receive(int).receive(int).send(bool)

It is all about duality!



receive(int).send(int).receive(bool)



receive(int).receive(int).send(bool)



Communication safety

No communication mismatch

Session Fidelity

• Communications follow the desired protocol

Progress

No deadlock/stuck in a session

Session Types and Objects 11) Mariangiola Shophia Drossopoulou NY 14-Dianitris h







Session Types and Objects

First Asynchronus Session

Shophia Drossopoulou

11)

2.

NY



TH

Diamitris



Session-based Distributed Programming in Java

Raymond Hu, Nobuko Yoshida

Kohei Honda

Imperial College London



Implementing Customer (4)

protocol p { begin. ![!<String>. ?(int) 1* ! { ACCEPT: { !<Address>. ?(Date) }, **REJECT:** { }

```
s.request();
s.outwhile(...) {
  s.send("PARIS/EUROSTAR");
  cost = s.receive();
}
if (...) {
  s.outbranch (ACCEPT) {
    s.send(...);
    date = s.receive();
  }
} else {
  s.outbranch(REJECT) { }
}
```

SJSocket s = **SJSocket.create(**p, ...);



CDL Equivalent

• Basic example:

package HelloWorld {

roleType YouRole, WorldRole; participantType You{YouRole}, World{WorldRole}; relationshipType YouWorldRel between YouRole and WorldRole; channelType WorldChannelType with roleType WorldRole;

```
choreography Main {
```

WorldChannelType worldChannel;

```
interaction operation=hello from=YouRole to=WorldRole
    relationship=YouWorldRel channel=worldChannel {
    request messageType=Hello;
}
```

Dr Gary Brown (Pi4 Tech) in 2007

Scribble Protocol

 "Scribbling is necessary for architects, either physical or computing, since all great ideas of architectural construction come from that unconscious moment, when you do not realise what it is, when there is no concrete shape, only a whisper which is not a whisper, an image which is not an image, somehow it starts to urge you in your mind, in so small a voice but how persistent it is, at that point you start scribbling" - Kohei Honda 2007

Basic example:

protocol HelloWorld { role You, World; Hello from You to World;





Session Types Overview



Global session type

$$G = A \rightarrow B : \langle U_1 \rangle. B \rightarrow C : \langle U_2 \rangle. C \rightarrow A : \langle U_3 \rangle$$

- Local session type
 - Slice of global protocol relevant to one role
 - Mechanically derived from a global protocol

$$T_A = !\langle B, U_1 \rangle.? \langle C, U_3 \rangle$$

- Process language
 - Execution model of I/O actions by session participants
 - Mechanically derived from a global protocol

 $P_A = a[A](x) \cdot x! \langle B, u_1 \rangle \cdot x?(C, y)$

• (Static) type checking for communication safety and progress



www.scribble.org

Home Getting Started Downloads Documentation - Community -

Scribble: Describing Multi Party Protocols

Scribble is a language to describe application-level protocols among communicating systems. A protocol represents an agreement on how participating systems interact with each other. Without a protocol, it is hard to do meaningful interaction: participants simply cannot communicate effectively, since they do not know when to expect the other parties to send data, or whether the other party is ready to receive data. However, having a description of a protocol has further benefits. It enables verification to ensure that the protocol can be implemented without resulting in unintended consequences, such as deadlocks.

Describe 🖋

Scribble is a language for describing multiparty protocols from a global, or endpoint neutral, perspective.

Verify 🖬

Scribble has a theoretical foundation, based on the Pi Calculus and Session Types, to ensure that protocols described using the language are sound, and do not suffer from deadlocks or livelocks.

Project 🗙

Endpoint projection is the term used for identifying the responsibility of a particular role (or endpoint) within a protocol.

Implement 🧮

Various options exist, including (a) using the endpoint projection for a role to generate a skeleton code, (b) using session type APIs to clearly describe the behaviour, and (c) statically verify the code against the projection.

Monitor **Q**

Use the endpoint projection for roles defined within a Scribble protocol, to monitor the activity of a particular endpoint, to ensure it correctly implements the expected behaviour.

Online tool : <u>http://scribble.doc.ic.ac.uk/</u>

```
module examples;
  1
  2
  3 - global protocol HelloWorld(role Me, role World) {
        hello() from Me to World;
  4
  5 -
       choice at World {
          goodMorning1() from World to Me;
  6
  7 -
       } or {
          goodMorning1() from World to Me;
  8
        }
  9
 10
     }
 11
                                             Role: Me
              Check Protocol: examples.HelloWorld
Load a sample 🔇
                                                                    Project
                                                                            Generate Graph
```

Multiparty Session Types (1)

POPL'08 Alice \rightarrow Bob: K (Nat) Bob \rightarrow Alice: K' (Int)

CONCUR'08 Alice -> Bob: (Nat) Bob -> Alice: (Int)



Multiparty Session Types (1) Much POPL'08 Simpler Alice -> Bob: K (Nat) Bob -> Alice: K (Int) CONCUR'08 30 Minutes Alice - Bob: (Nat) Bob -> Alice: (Int) Connection

Multiparty Session Types (1) Much POPL'08 Simpler Alice -> Bob: K (Nat) Bob -> Alice: K (Int) CONCUR'08 30 Minutes My 4th 80+ Alice -> Bob: (Nat) Bob -> Alice: (Int) Connection

CONCUR'08 Global Progress of MPST



COORDINATION'13 Inference of Grlobal Progress in MPST

MSCS Gentle Introduction of MPST



$Us \in \mathbf{M}obility \, \mathbf{R}esearch \, \mathbf{G}roup$

	tyReadingGroup Session Types research at Imperial College	
Home People Public	ations Grants Talks Tools Awards Kohei Honda	
NEWS	SELECTED	
Our recent work Fencing off Go: Liveness and Safety for Channel- based Programming was summarised on The Morning Paper blog.	PUBLICATIONS	
2 Feb 2017 Weizhen passed her viva today,	Raymond Hu , Nobuko Yoshida : Explicit Connection Actions in Multiparty Session Types. <i>To appear in</i> FASE 2017 .	
congratulations Dr. Yang! 24 Jan 2017	Julien Lange , Nicholas Ng , Bernardo Toninho , Nobuko Yoshida : Fencing off Go: Liveness and Safety for Channel-based Programming. POPL 2017 .	
Mariangiola Dezani-Ciancaglini, a long-term collaborator with our group working on Session Types turns 70 today, more details here.	Rumyana Neykova , Nobuko Yoshida : Let It Recover: Multiparty Protocol- Induced Recovery. CC 2017 .	
23 Dec 2016 Rumyana passed her viva today,	Julien Lange , Nobuko Yoshida : On the Undecidability of Asynchronous Session Subtyping. <i>To appear in</i> FoSSaCS 2017 .	

http://mrg.doc.ic.ac.uk/

Academic Staff Nobuko Yoshida **Research Associate** Raymond Hu Julien Lange Nicholas Ng Xinyu Niu **Alceste Scalas** Bernardo Toninho PhD Student Assel Altayeva Juliana Franco

Rumyana Neykova

Weizhen Yang

Ocean Observatories Initiative

- A NSF project (400M\$, 5 Years) to build a cyberinfrastructure for observing oceans around US and beyond.
- Real-time sensor data constantly coming from both off-shore and on-shore (e.g. buoys, submarines, under-water cameras, satellites), transmitted via high-speed networks.





Ocean Observatories Initiative





OOI agent negotiation 1/5



https://confluence.oceanobservatories.org/display/syseng/ CIAD+COI+OV+Negotiate+Protocol

・ロト <
ゆ ト <
き ト <
き ト き
の へ で
11/42
</p>

OOI agent negotiation 2/5

type <yml> "SAPDoc1" from "SAPDoc1.yml" as SAP;

global protocol Negotiate(role Consumer as C, role Producer as P) {



OOI agent negotiation 3/5 (choice)

type <yml> "SAPDoc1" from "SAPDoc1.yml" as SAP;

```
global protocol Negotiate(role Consumer as C, role Producer as P) {
   propose(SAP) from C to P;
```

```
choice at P {
   accept() from P to C;
   confirm() from C to P;
} or {
   reject() from P to C;
} or {
   propose(SAP) from P to C;
```


OOI agent negotiation 4/5

type <yml> "SAPDoc1" from "SAPDoc1.yml" as SAP;

```
global protocol Negotiate(role Consumer as C, role Producer as P) {
   propose(SAP) from C to P;
```

```
choice at P {
 accept() from P to C;
 confirm() from C to P;
r f
 reject() from P to C;
} or {
 propose(SAP) from P to C;
 choice at C {
   accept() from C to P;
   confirm() from P to C;
 f or f
   reject() from C to P;
 f or f
   propose(SAP) from C to P;
```



OOI agent negotiation 5/5 (recursion)

type <yml> "SAPDoc1" from "SAPDoc1.yml" as SAP;

```
global protocol Negotiate(role Consumer as C, role Producer as P) {
    propose(SAP) from C to P;
    rec X {
         choice at P {
             accept() from P to C;
                                                                                                             Consumer
                                                                                                                                                  Provider
             confirm() from C to P;
                                                                                                               Agent
                                                                                                                                                   Agent
                                                                                                                           negotiate: request(SAP 1)
                                                                                                                                                              Confirm is the
                                                                                                Negotiation starting by a
        r f
                                                                                                                                                           complementary accept
                                                                                              Consumer making a proposal, then accepted by Provider and
                                                                                                                         negotiate: accept(SAP 1, details)
                                                                                                                                                           by the other party (both
                                                                                                 nfirmed by Consume
                                                                                                                                                            must accept for an
                                                                                                                           negotiate: confirm(SAP 1)
                                                                                                                                                              agreement)
            reject() from P to C;
                                                                                                                                                          With a mutual accept, a
        } or {
                                                                                                                                                           least one commitment
                                                                                                             ALT
                                                                                                                                                           on each side of the
                                                                                                                            negotiate: invite(SAP 1)
                                                                                                                                                            conversation results
                                                                                               Negotiation starting by the
                                                                                                                                                           (may be multiple). The
            propose(SAP) from P to C;
                                                                                                                         negotiate: accept(SAP 1, details)
                                                                                                 ider inviting a Consumer
                                                                                                                                                           contract is as stated in
                                                                                               with a proposal, accepted by 
Consumer and confirmed by
                                                                                                                           negotiate: confirm(SAP 1)
                                                                                                                                                           the most recent SAP
             choice at C {
                                                                                                             ALT
                                                                                                                           negotiate: request(SAP_1)
                 accept() from C to P;
                                                                                                                                                           A counter-propose is a
                                                                                                                                                           new SAP, but it typically
                                                                                                Negotiation starting by a
                                                                                                                        negotiate: counter-propose(SAP_2)
                                                                                                                                                            refines or partially
                                                                                              Consumer making a proposal.
                  confirm() from P to C;
                                                                                                                                                           modifies the prior SAP
                                                                                              The recipient (Provider) makes
                                                                                                                         negotiate: accept(SAP 2, details)
                                                                                              a counter-proposal, supplanting
                                                                                              SAP 1, which is then accepted
                                                                                                                           negotiate: confirm(SAP_2)
                                                                                              by Consumer and confirmed by
                                                                                                  the Provider.
             \mathbf{r}
                 reject() from C to P;
                                                                                                             ALT
                                                                                                                           negotiate: request(SAP 1)
                                                                                                                                                            Any party can reject
                                                                                                                                                            instead of counter-
                                                                                                Negotiation starting by a
                                                                                                                            negotiate: reject(SAP_1)
                                                                                              Consumer making a proposal,
                                                                                                                                                           propose (or accept)
                                                                                              rejected by the Provider ending
             f or f
                                                                                                  the Neootiation
                 propose(SAP) from C to P;
                  continue X;
             }
                                                                                                             ◆□▶ ◆□▶ ◆目▶ ◆目▶ ● ● のへで
         }
                                                                                                                                                                    15/42
```

Local protocol projection (Negotiation Consumer)

```
// Global
propose(SAP) from C to P;
rec START {
 choice at P {
   accept() from P to C;
   confirm() from C to P;
 \mathbf{r}
   reject() from P to C;
 f or f
   propose(SAP) from P to C;
   choice at C {
     accept() from C to P;
     confirm() from P to C;
   \mathbf{r}
     reject() from C to P;
   } or {
     propose(SAP) from C to P;
     continue START;
```

// Projection for Consumer propose(SAP) to P; rec START { choice at P { accept() from P; confirm() to P; } or { reject() from P; } or { propose(SAP) from P; choice at C { accept() to P; confirm() from P; } or { reject() to P; } or { propose(SAP) to P; continue START;

FSM generation (Negotiation Consumer)



▲□▶ ▲□▶ ▲ ■▶ ▲ ■ シ ● ② へ (~ 20 / 42)



RESOURCES

- All Resources
- Data Products
- 88 Observatories
- Platforms
- Instruments

Welcome to Release 2 of the Ocean Observatories Initiative Observatory (OOI). You already have access to many OOI features and real-time data. Just click on something that looks interesting on this page to start using the OOI as our Guest.

For personalized services, such as setting up notifications and preserving settings for your next visit, create a free account by clicking on "Create Account" at the top of the page.



National Science Foundation working with Consortium for Ocean Leadership

Funding for the Ocean Observatories Initiative is provided by the National Science Foundation through a Cooperative Agreement with the Consortium for Ocean Leadership. The OOI Program Implementing Organizations are funded through sub-awards from the Consortium for Ocean Leadership.

Location CURRENT LOCATION

O Temperature

O Salinity

O Oxygen

O Density

O Currents

O Turbidity

O pH

O Other

O Chlorophyll

O Seismology

Sea Surface Height (SSH)



RECENT UPDATES

RECENCY

1 Hour

2 hours

3 hours

5 hours

8 hours

12 hours

18 hours

24 hours

48 Hours

72 Hours

1

1

~

	NAME	DATE	TYPE	EVENT	DESCRIPTION	NOTE
O 01 m	Oregon Coast North Salinity	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
0 01 m	California South 100m pH	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
0 01 m	California South salinity	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
O 03 m	Oregon North Turbidity	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
O 05 m	Oregon SouthTemperature	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
O 20 m	Oregon Coast Currents	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
0 01 h	California South Seismology	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
O 01 h	Oregon Coast South 1000m Ox	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
02 h	California Coast Seismology	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here
0 04 h	California North Seismology	2012-01-10 23:55:55	Type	Event	Description goes here	Note goes here

Dashboard

RECENT IMAGES



Last Modified: 2011-06-15 Last Viewed: 2011-12-15 Last Updated: 2011-12-30, 13.24



Gorgonian Coral Last Modified: 2011-06-15 Last Viewed: 2011-12-15 Last Updated: 2011-12-30, 13.24

Acoustic Release

Last Modified: 2011-06-15



POPULAR RESOURCES

SeaBird CDT

Last Modified: 2011-06-15 Last Viewed: 2011-12-15 Last Updated: 2011-12-30, 13.24

Marine caption Last Modified: 2011-06-15 Last Viewed: 2011-12-15 Last Updated: 2011-12-30, 13.24



Surface Buoy Last Modified: 2011-06-15 Last Viewed: 2011-12-15 Last Updated: 2011-12-30, 13.24

UNUSUAL EVENTS



Oregon Coast Wave Heigh Last Modified: 2011-06-15 Last Viewed: 2011-12-15 Last Updated: 2011-12-30, 13.24



Water Surface Elevation Last Modified: 2011-06-15 Last Viewed: 2011-12-15 Last Updated: 2011-12-30, 13.24

FACEPAGE

COMPOSITE



Dynamic Monitoring [RV'13, COORDINATION'14, FMSD'15, LMCS'17, CC'17]



Type Checking [OOPSLA'15, ECOOP'16, ECOOP'17, COORDINATION'17]



Code Generation [CC'15, FASE'16, FASE'17]



Synthesis [ICALP'13, POPL'15, CONCUR'15, TACAS'16, CC'16]





Session Type based Tools

OOI Governance



ZDLC: Process Modeling



Actor Verification



MPI code generations



Interactions with Industries







Adam Bowen @adamnbowen · Sep 15 I didn't even know that session types existed an hour ago, but thanks to Nobuko Yoshida's great talk at **#pwlconf**, I want to learn more.

Nobuko Yoshida Imperial College, London

DoC researcher to speak at Golang UK conference

by Vicky Kapogianni 20 July 2016



DoC researcher to speak at industry-focused Golang UK conference on results of concurrency research

Click here to add content

 @nicholascwng rocking on @GolangUKconf about static deadlock detection in #golang #gouk16



Interactions with Industries

F#unctional Londoners Meetup Group

6 days ago · 6:30 PM Session Types with Fahd Abdeljallal



43 Members

Synopsis: Session types are a formalism to codify the structure of a communication, using types to specify the communication protocol used. This formalism provides the... LEARN MORE

Distributed Systems vs.

Compositionality

Dr. Roland Kuhn @rolandkuhn — *CTO of Actyx*

actyx

Current State

- behaviors can be composed both sequentially and concurrently
- effects are not yet tracked
- Scribble generator for Scala not yet there
- theoretical work at Imperial College, London (Prof. Nobuko Yoshida & Alceste Scalas)





Go concurrency verification research at DoC grabs headline

A paper by DoC researchers at POPL on Go concurrency verification was featured in a tech blog and generates a buzz outside of the research community.

A paper by researchers at the department was recently featured in the morning paper, a blog by venture capitalist Adrian Colye, which summarises an important, influential, topical or otherwise interesting paper in the field of computer science every weekday in an easily digestible way by non-researchers. On the 2 Feb 2017 issue of the morning paper, It was highlighted as "the true spirit of POPL (Principles of Programming Languages)".

Java API Generation [FASE'16] \mathcal{T}

RFC 821 Sim	August 1982 ple Mail Transfer Protocol							
TABLE OF CONTENTS								
1. INTRODUCTION	<u>1</u>							
2. THE SHTP HODEL								
3. THE SMTP PROCEDURE	4							
3.1. Mail 3.2. Forwarding 3.3. Verifying and Expanding 3.4. Sending and Mailing 3.5. Opening and Closing 3.6. Relaying 3.7. Domains 3.8. Changing Roles 4. THE SMTP SPECIFICATIONS 4.1. SMTP Commands 4.1.1. Command Semantics 4.1.2. Command Syntax 4.2. SMTP Replies 4.2.1. Reply Codes by Function Group 4.2.2. Reply Codes in Numeric Order 4.3. Sequencing of Commands and Replies 4.4. State Diagrams 4.5. Details 4.5. Transparency 4.5.3. Sizes	7 8 11 13 14 17 18 19 11 13 14 15 36 37 39 41 41							



Scribble - Proving a distributed design





× Information Flow and Security [CONCUR'10]....

× Reversible Computation ECONCUR'17]...

× Preciseness EPPPP/14] ELMCSJ...

Types and Subtypes for Client-Server Interactions

Simon Gay and Malcolm Hole

(ESOP'99)

<ロ> (四) (四) (三) (三) (三) (三)



(中) (종) (종) (종) (종) (종)





Other completeness results Subtyping of recursive types requires algorithms for checking subtype relations, as discussed in [32, Chapter 21]. These algorithms need to be proved sound and complete with respect to the definition of the corresponding subtyping, as done for example in [7, 12, 33]. Algorithms for checking the synchronous and asynchronous subtypings of the present paper can be easily designed.



× Information Flow and Security [CONCUR'10]....

× Reversible Computation ECONCUR'17]...

× Preciseness EPPPP/14] ELMCSJ...

